# CSCE 5350.001 Fundamentals of Database Systems Project Part 5

Naga Vara Pradeep Yendluri

11646461

nagavarapradeepyendluri@my.unt.edu

# **Project Description:**

The Movie Producer Management System is an application that is being developed for a movie production company like Universal Studios. The system is designed to store and manage information about the company's movies, artists, songs, employees, and various other aspects of the movie production process. The system will store information about the producing site locations, movie-script-inventory, sponsoring companies, employee data, and payroll. It will also store information about the artists and the movies they have worked on, as well as the various aspects of the movie production process, such as soundtracks, awards, and more.

We have Identified the following entities and relations for the movie producer management system. We have identified the Functional Dependencies to make the database normalized to BCNF or 3NF and identified the dependency preserving and lossless join between the tables also updated the tables accordingly.

- 1. **Movies:** The entity 'Movies' provides information about the different movies produced by the company. It has 6 attributes, including the movie id, movie title, release date, duration, in production and director. This entity is important for keeping track of the different movies produced by the company and the information related to each movie.
- 2. Artists: The entity 'Artists' provides information about the actors involved in the movies. It has 5 attributes, including the actors id, actors name, actors date of birth, address and age. This entity is important for maintaining the information about the actors, their age and date of birth, which is required for casting actors for various roles.
- 3. **Genre:** The entity 'Genre' provides information about the genre to which the movie belongs. It has 2 attributes, genre id, and genre name. This entity is important for categorizing the movies into different genres, which helps in better management and analysis of the movies.
- 4. **Sponsoring Companies:** The entity 'Sponsoring Companies' provides information about the companies that sponsor the movies. It has 4 attributes, including sponsor id, sponsor name, movie id and movie sponsored. This entity is important for tracking the sponsorship deals and the companies that sponsor the movies.

- 5. **Site Locations:** The entity 'Site Locations' provides information about the different producing sites, including their addresses and buildings. It has 4 attributes, including location id, name, address, and building names. This entity is important for tracking the different producing sites and their details, which is essential for managing the movie production process.
- 6. **Buildings:** The entity 'Buildings' provides information about the buildings in each producing site. It has 3 attributes, including Building id, name, and type of building. This entity is important for tracking the different types of buildings present in the producing sites, which is essential for managing the resources and maintenance of the buildings.
- 7. **Movie Script Inventory:** The entity 'Movie Script Inventory' provides information about the movie scripts. It has 5 attributes, including script id, name, movie id, author, and publication date. This entity is important for tracking the different movie scripts and the information related to each script.
- 8. **Employee:** The entity 'Employee' provides information about the employees of the company. It has 5 attributes, including employee id, name, job title, hourly pay and phone. This entity is important for maintaining the information about the employees, their job title and contact information, which is required for managing the human resources of the company.
- 9. **Payroll:** The entity 'Payroll' provides information about employee payroll data. It has 5 attributes, including employee id, hours worked, joining date, work date, etc. This entity is important for tracking the payroll information of the employees, which is essential for managing the finances of the company.
- 10. **Songs:** The entity 'Songs' provides information about different soundtracks used in the movies. It has 4 attributes, including song id, track title, movie id, and singer name. This entity is important for tracking the different soundtracks used in the movies and the information related to each soundtrack.

# Latest Schema:

Movie (movie\_id, title, avg\_rating, date\_of\_release, duration, script\_inventory\_id, director, in\_production)

Songs (song\_id, song\_name, singer\_name, movie\_id)

Genre (genre\_id, genre\_name)

TaggedWith (movie\_id, genre\_id)

SiteLocation (location\_id, location\_name, address)

ShotAt (location\_id, movie\_id)

Building (building\_id, building\_name, purpose, location\_id)

PostProductionDoneIn (movie\_id, building\_id)

Employees (employee\_id, employee\_name, designation, phone\_number, hourly\_pay, employee\_level)

Manages (employee\_id, location\_id) Payroll (payroll\_id, employee\_id, hours\_worked, date) SponsoringCompany (company\_id, company\_name, num\_movies\_produced) getsPaidBy (artist\_id, company\_id) Produces (company\_id, movie\_id) Artist (artist\_id, artist\_name, date\_of\_birth, gender, address) MovieScriptInventory (script\_inventory\_id, script\_inventory\_name) ActsIn(movie\_id, artist\_id) Rating(rating id, movie id, rating)

## **Additional Assumptions:**

- From the latest schema we got 18 functional dependencies.
- Normalized the tables into 2NF,3NF and BCNF.

#### Functional dependencies for the database:

Movie: movie\_id  $\rightarrow$  title, date\_of\_release, duration, script\_inventory\_id, director, in\_production MovieRating: movie\_id  $\rightarrow$  avg\_rating Songs: song\_id  $\rightarrow$  song\_name, singer\_name, movie\_id Genre: genre\_id  $\rightarrow$  genre\_name MovieGenre: (movie\_id, genre\_id)  $\rightarrow$  {movie\_id, genre\_id} SiteLocation: location\_id  $\rightarrow$  location\_name, address MovieLocation: (location\_id, movie\_id)  $\rightarrow$  {location\_id, movie\_id} Building: building\_id  $\rightarrow$  building\_name, purpose, location\_id MoviePostProduction: (movie\_id, building\_id)  $\rightarrow$  {movie\_id, building\_id} Employees: employee\_id  $\rightarrow$  employee\_name, designation, phone\_number, employee\_level EmployeeSalary: employee id  $\rightarrow$  hourly pay Manages: (employee\_id, location\_id) → {employee\_id, location\_id} Payroll: payroll\_id → employee\_id, hours\_worked, date, hourly\_pay SponsoringCompany: company\_id → company\_name, num\_movies\_produced Artist: artist\_id → artist\_name, date\_of\_birth, gender, address ArtistCompany: (artist\_id, company\_id) → {artist\_id, company\_id} MovieScriptInventory: script\_inventory\_id → script\_inventory\_name ActsIn: (movie\_id, artist\_id) → {movie\_id, artist\_id} From the above functional dependencies, we are taking 12 among them.

# Normalization:

• <u>1NF:</u>

The tables are in 1NF since they do not have repeating groups or multivalued attributes.

• <u>2NF:</u>

## Movie:

Movie: movie\_id → title, date\_of\_release, duration, script\_inventory\_id, director, in\_production

Candidate key, prime attributes = { movie\_id }

Non – prime attributes = { in\_production }

Since, both the location\_name and address are fully dependent on location\_id. There are no partial dependencies, hence it is in 2NF.

## Songs:

 $song_id \rightarrow song_name$ ,  $singer_name$ ,  $movie_id$ 

Candidate key, prime attributes = { song\_id }

Non – prime attributes = { song\_name, singer\_name, movie\_id }

All the non – prime attributes are fully dependent on candidate key (song\_id). So, it is in 2NF

## Genre:

 $genre\_id \rightarrow genre\_name$ 

Candidate key, prime attributes = { genre\_id }

Non – prime attributes = { genre\_name }

Since, the FD consists of only one attribute on both sides, so it is trivially in 2NF.

## Site location:

location\_id  $\rightarrow$  location\_name, address

Candidate key, prime attributes = { location\_id }

Non – prime attributes = { location\_name, address }

All the non-prime attributes are fully dependent on the candidate key this is in 2NF.

## Payroll:

payroll\_id  $\rightarrow$  employee\_id, hours\_worked, date, hourly\_pay

Candidate key, prime attributes = { payroll\_id }

Non – prime attributes = { employee\_id, hours\_worked, date, hourly\_pay }

All the non-prime attributes are fully dependent on primary key it is in 2NF.

## Artist:

artist\_id  $\rightarrow$  artist\_name, date\_of\_birth, gender, address

Candidate key, prime attributes = { artist\_id }

Non – prime attributes = { artist\_name, date\_of\_birth, gender, address }

All the non-prime attributes are fully dependent on the candidate key. So, it is in 2NF.

## ActsIn:

movie\_id, artist\_id  $\rightarrow$  movie\_id, artist\_id

Candidate key, prime attributes = { artist\_id, movie\_id }

Non – prime attributes = None

Hence there are no partial dependencies and non-prime attributes it is in 2NF.

## MoviePostProduction:

movie\_id, building\_id → movie\_id, building\_id

Candidate key, prime attributes = { movie\_id, building\_id }

Non – prime attributes =None

Hence there are no partial dependencies and non-prime attributes it is in 2NF.

## **Employees:**

 $employee\_id \rightarrow employee\_name, designation, phone\_number, employee\_level$ 

Candidate key, prime attributes = { employee\_id }

Non – prime attributes ={ employee\_name, designation, phone\_number, employee\_level }

As there is only one candidate key and no partial dependencies on any part of the candidate key, this relation is already in 2NF.

#### MovieScriptInventory:

 $script\_inventory\_id \rightarrow script\_inventory\_name$ 

Candidate key, prime attributes ={ script\_inventory\_id }

Non – prime attributes =None

Since the given FD has only one attribute on the right-hand side, it is already in 2NF and no further normalization is needed.

## • <u>3NF:</u>

#### Movie:

Movie: movie\_id o title, date\_of\_release, duration, script\_inventory\_id, director, in\_production

Candidate key, prime attributes = { movie\_id }

```
Non – prime attributes = { in_production }
```

**Movie\_id** is the **super key**. There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### Songs:

```
song id \rightarrow song name, singer_name, movie_id
```

Candidate key, prime attributes = { song\_id }

Non – prime attributes = { song\_name, singer\_name, movie\_id }

Any combination of attributes that includes **song\_id** would be a **super key.** 

#### Genre:

genre\_id  $\rightarrow$  genre\_name

Candidate key, prime attributes = { genre\_id }

Non – prime attributes = { genre\_name }

**Genre\_id** is the **super key**. There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### Site location:

location\_id  $\rightarrow$  location\_name, address

Candidate key, prime attributes = { location\_id }

Non – prime attributes = { location\_name, address }

**Location\_id** is the **super key.** There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### **Payroll:**

payroll\_id  $\rightarrow$  employee\_id, hours\_worked, date, hourly\_pay

Candidate key, prime attributes = { payroll\_id }

Non – prime attributes = { employee\_id, hours\_worked, date, hourly\_pay }

Any combination of attributes that includes **payroll\_id** would be a **super key**.Since there are no transitive dependencies. Hence, the given relation is in 3NF.

#### Artist:

artist\_id  $\rightarrow$  artist\_name, date\_of\_birth, gender, address

Candidate key, prime attributes = { artist\_id }

Non – prime attributes = { artist\_name, date\_of\_birth, gender, address }

Any combination of attributes that includes **artist\_id** would be a **super key**. There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### ActsIn:

movie\_id, artist\_id  $\rightarrow$  movie\_id, artist\_id

Candidate key, prime attributes = { artist\_id, movie\_id }

Non – prime attributes = None

Any combination of attributes that includes both **artist\_id** and **movie\_id** would be a **superkey** for this relation. There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### **MoviePostProduction:**

movie\_id, building\_id  $\rightarrow$  movie\_id, building\_id

Candidate key, prime attributes = { movie\_id, building\_id }

Non – prime attributes =None

The super key is { movie\_id, building\_id }.

There are no transitive dependencies in the relation. Hence, the given relation is in 3NF.

#### **Employees:**

 $employee_id \rightarrow employee_name$ , designation, phone\_number, employee\_level

Candidate key, prime attributes = { employee\_id }

Non – prime attributes ={ employee\_name, designation, phone\_number, employee\_level }

This functional dependency is already in 3NF as there are no transitive dependencies.

#### MovieScriptInventory:

script\_inventory\_id → script\_inventory\_name

Candidate key, prime attributes ={ script\_inventory\_id }

Non – prime attributes = { }

Since there is only one FD it will satisfy the 3NF condition and it is in 3NF.

## • <u>BCNF:</u>

#### Movie:

Movie: movie\_id → title, date\_of\_release, duration, script\_inventory\_id, director, in\_production

Candidate key, prime attributes = { movie\_id }

Non – prime attributes = { in\_production }

The only candidate key is { movie\_id }, and there are no non-trivial dependencies. Therefore, the relation is in BCNF.

#### Songs:

 $song_id \rightarrow song_name$ ,  $singer_name$ ,  $movie_id$ 

Candidate key, prime attributes = { song\_id }

Non – prime attributes = { song\_name, singer\_name, movie\_id }

The only candidate key is { song\_id }, and there are no non-trivial dependencies. Therefore, the relation is in BCNF.

#### Genre:

genre\_id  $\rightarrow$  genre\_name

Candidate key, prime attributes = { genre\_id }

Non – prime attributes = { genre\_name }

The only candidate key is { genre\_id }, and there are no non-trivial dependencies. Therefore, the relation is in BCNF.

#### Site location:

location\_id  $\rightarrow$  location\_name, address

Candidate key, prime attributes = { location\_id }

Non – prime attributes = { location\_name, address }

The only candidate key is { location\_id }, and there are no non-trivial dependencies. Therefore, the relation is in BCNF.

## **Payroll:**

payroll\_id  $\rightarrow$  employee\_id, hours\_worked, date, hourly\_pay

Candidate key, prime attributes = { payroll\_id }

Non – prime attributes = { employee\_id, hours\_worked, date, hourly\_pay }

The given relation is not in BCNF. To bring it to BCNF, we need to decompose the relation into two relations:

- R1(employee\_id, hourly\_pay)
- R2(payroll\_id, employee\_id, hours\_worked, date)

Each relation has a single determinant for each of its attributes, and the relation satisfies BCNF.

## Artist:

artist\_id  $\rightarrow$  artist\_name, date\_of\_birth, gender, address

Candidate key, prime attributes = { artist\_id }

Non – prime attributes = { artist\_name, date\_of\_birth, gender, address }

Since there is only one candidate key and no functional dependencies other than the trivial ones, the relation is automatically in BCNF.

## ActsIn:

movie\_id, artist\_id  $\rightarrow$  movie\_id, artist\_id

Candidate key, prime attributes = { artist\_id, movie\_id }

Non – prime attributes = None

Since there are no non-prime attributes and every non-trivial functional dependency in the relation has a candidate key as the determinant, the given relation is in BCNF.

## MoviePostProduction:

movie\_id, building\_id  $\rightarrow$  movie\_id, building\_id

Candidate key, prime attributes = { movie\_id, building\_id }

Non – prime attributes =None

Since there are no prime attributes, it is in BCNF.

## **Employees:**

employee\_id → employee\_name, designation, phone\_number, employee\_level

Candidate key, prime attributes = { employee\_id }

Non – prime attributes ={ employee\_name, designation, phone\_number, employee\_level }

It is also in BCNF since the left-hand side (LHS) contains the candidate key.

## MovieScriptInventory:

script\_inventory\_id → script\_inventory\_name

Candidate key, prime attributes ={ script\_inventory\_id }

Non – prime attributes ={ }

As the given FD has only one candidate key it is the super key, so it is in BCNF.

Therefore, all the Functional dependencies are in BCNF.

# **Dependency Preserving and Lossless Join:**

## Movie:

Movie: movie\_id  $\rightarrow$  title, date\_of\_release, duration, script\_inventory\_id, director, in\_production Here we have Movie table in which we have attributes of movie\_script\_invetory table. Comparing both tables to find whether they have common attributes and dependencies. Movie:

- movie\_id  $\rightarrow$  title
- movie\_id  $\rightarrow$  date\_of\_release
- movie\_id  $\rightarrow$  duration
- movie\_id  $\rightarrow$  script\_inventory\_id
- movie\_id  $\rightarrow$  director

• movie\_id  $\rightarrow$  in\_production

Script\_Inventory:

• script\_inventory\_id  $\rightarrow$  {}

The common attribute between the two tables is "script\_inventory\_id". However, since the "script\_inventory\_id" attribute is a foreign key in the "Movie" table.

Therefore, based on the information given, we can conclude that the join between the "Movie" and "Script\_Inventory" tables is both dependency preserving and lossless.

## Songs:

```
song_id \rightarrow song_name, singer_name, movie_id
```

If the "movie\_id" attribute in the "song" table refers to a separate table containing information about the movies that the songs belong to, we have the following dependencies in each table:

Song:

- $song_id \rightarrow song_name$
- $song_id \rightarrow singer_name$
- $song_id \rightarrow movie_id$

Movie:

• movie\_id  $\rightarrow$  {}

The common attribute between the two tables is "movie\_id".

For example, if a movie has two songs with different song\_ids but the same singer and song name, then the join would produce an additional tuple that combines the information for those two songs.

Therefore, we can conclude that the join between the "Song" and "Movie" tables is dependency preserving and lossless join.

#### Genre:

• genre\_id  $\rightarrow$  genre\_name

There are no other tables to join with, so there are no common attributes to consider.

Based on the above scenario this is dependency preserving,

Since the dependency is a simple one-to-one mapping between genre\_id and genre\_name, there are no redundant tuples in the table that can be eliminated.

Therefore, we can conclude that the table is lossless.

#### Site location:

location\_id  $\rightarrow$  location\_name, address

Location:

- location\_id  $\rightarrow$  location\_name
- location\_id  $\rightarrow$  address

There are no other tables to join with, so there are no common attributes to consider.Based on these dependencies, we can conclude that the table is dependency preserving, as the dependencies are trivially preserved by any join operation.

To determine if the table is lossless, we need to check if the table if there are any redundant tuples, as there are no redundant tuples the join is lossless.

## **Payroll:**

 $payroll_id \rightarrow employee_id$ , hours\_worked, date, hourly\_pay

These are the individual dependencies.

Payroll:

- payroll\_id  $\rightarrow$  employee\_id
- payroll\_id  $\rightarrow$  hours\_worked
- payroll\_id  $\rightarrow$  date
- $payroll_id \rightarrow hourly_pay$

From the above dependencies we conclude that it is dependency preserving.

Since payroll\_id is the primary key of the table, there are no redundant tuples in the table that can be eliminated without losing any information. Therefore, we can conclude that the table is lossless.

The "Payroll" table is both dependency preserving and lossless.

## Artist:

artist\_id  $\rightarrow$  artist\_name, date\_of\_birth, gender, address

The only table we have is the one containing the information about artists. In this case, we can identify the following dependency in the table:

Artists:

- $artist_id \rightarrow artist_name$
- $artist_id \rightarrow date_of_birth$

- $artist_id \rightarrow gender$
- $artist_id \rightarrow address$

Based on these dependencies, we can conclude that the table is dependency preserving, as all of the dependencies are preserved by any join operation.

The "Artists" table is lossless, since there are no redundant tuples that can be eliminated without losing any information.

#### ActsIn:

movie\_id, artist\_id  $\rightarrow$  movie\_id, artist\_id

The combination of values in movie\_id and artist\_id uniquely determines the values of movie\_id and artist\_id. This dependency is trivially true, as the right-hand side of the dependency is simply equal to the left-hand side.

The join of the two tables on the common attribute movie\_id and artist\_id will result in a table with the same attributes as the original tables.

Therefore, the join between these two tables is both dependency preserving and lossless.

#### **MoviePostProduction:**

movie\_id, building\_id → movie\_id, building\_id

The movie\_id and building\_id uniquely determine the values of movie\_id and building\_id. This dependency is trivially true.

Both the primary keys movie\_id and building \_id, the join will result in the same table without any information loss.

Hence, the MoviePostProduction is lossless and dependency preserving.

#### **Employees:**

employee\_id → employee\_name, designation, phone\_number, employee\_level

Employees:

- employee\_id  $\rightarrow$  employee\_name
- employee\_id  $\rightarrow$  designation
- employee\_id  $\rightarrow$  phone\_number
- $employee_id \rightarrow employee_level$

Based on these dependencies, we can conclude that the table is dependency preserving, as all of the dependencies are preserved by any join operation.

The "Employees" table is lossless, since there are no redundant tuples that can be eliminated without losing any information.

## MovieScriptInventory:

```
script_inventory_id → script_inventory_name
```

Since we have no other table, the dependency is preserved.

Since the script\_inventory\_id is a primary key, there can be no duplicate values of script\_inventory\_id in the table. Therefore, there can be no redundant tuples in the table, and the table is lossless.

## New tables:

#### Movie:

CREATE TABLE Movie (

movie\_id INT PRIMARY KEY,

title VARCHAR2(100) NOT NULL,

date\_of\_release DATE NOT NULL,

duration INT NOT NULL,

script\_inventory\_id INT REFERENCES MovieScriptInventory(script\_inventory\_id) NOT
NULL,

director VARCHAR2(100) NOT NULL,

```
in_production CHAR(1) NOT NULL
```

);

## **MovieRating:**

CREATE TABLE MovieRating (

movie\_id INT PRIMARY KEY REFERENCES Movie(movie\_id),

avg\_rating NUMBER(3, 1) NOT NULL

);

## Songs:

CREATE TABLE Songs (

song\_id INT PRIMARY KEY,

song\_name VARCHAR2(100) NOT NULL, singer\_name VARCHAR2(100) NOT NULL, movie\_id INT REFERENCES Movie(movie\_id) NOT NULL );

#### Genre:

CREATE TABLE Genre (

genre\_id INT PRIMARY KEY,

genre\_name VARCHAR2(100) NOT NULL

);

#### MovieGenre:

CREATE TABLE MovieGenre ( movie\_id INT REFERENCES Movie(movie\_id), genre\_id INT REFERENCES Genre(genre\_id), PRIMARY KEY(movie\_id, genre\_id)

);

#### SiteLocation:

CREATE TABLE SiteLocation ( location\_id INT PRIMARY KEY, location\_name VARCHAR2(100) NOT NULL, address VARCHAR2(200) NOT NULL );

#### **MovieLocation:**

CREATE TABLE MovieLocation ( location\_id INT REFERENCES SiteLocation(location\_id), movie\_id INT REFERENCES Movie(movie\_id), PRIMARY KEY(location\_id,movie\_id) **Building:** 

CREATE TABLE Building ( building\_id INT PRIMARY KEY, building\_name VARCHAR2(100) NOT NULL );

## **Employees:**

CREATE TABLE Employees ( employee\_id INT PRIMARY KEY, hourly\_pay NUMBER(6, 2) NOT NULL, UNIQUE(phone\_number) );

## **EmployeeDetails:**

CREATE TABLE EmployeeDetails ( employee\_id INT REFERENCES Employees(employee\_id), employee\_name VARCHAR2(100) NOT NULL, designation VARCHAR2(100) NOT NULL, employee\_level VARCHAR2(100) NOT NULL, PRIMARY KEY(employee\_id) );

## Manages:

CREATE TABLE Manages ( employee\_id INT REFERENCES Employees(employee\_id), location\_id INT REFERENCES SiteLocation(location\_id),

);

#### PRIMARY KEY(employee\_id, location\_id)

);

## **Payroll:**

CREATE TABLE Payroll ( payroll\_id INT PRIMARY KEY, employee\_id INT REFERENCES Employees(employee\_id) NOT NULL, hours\_worked NUMBER(3, 1) NOT NULL, date DATE NOT NULL );

#### **Sponsoring\_Company:**

CREATE TABLE SponsoringCompany ( company\_id INT PRIMARY KEY, company\_name VARCHAR2(100) NOT NULL );

## **CompanyMovies:**

CREATE TABLE CompanyMovies ( company\_id INT REFERENCES SponsoringCompany(company\_id), movie\_id INT REFERENCES Movie(movie\_id), PRIMARY KEY(company\_id, movie\_id) );

#### getsPaidBy:

CREATE TABLE getsPaidBy ( artist\_id INT PRIMARY KEY REFERENCES Artist(artist\_id), company\_id INT REFERENCES SponsoringCompany(company\_id) NOT NULL

#### **Produces:**

CREATE TABLE Produces ( company\_id INT REFERENCES SponsoringCompany(company\_id), movie\_id INT REFERENCES Movie(movie\_id), PRIMARY KEY(company\_id, movie\_id) );

## ArtistDetails:

CREATE TABLE ArtistDetails ( artist\_id INT REFERENCES Artist(artist\_id), date\_of\_birth DATE NOT NULL, gender CHAR(1) NOT NULL, address VARCHAR2(200) NOT NULL, PRIMARY KEY(artist\_id) );

## ActsIn:

CREATE TABLE ActsIn ( movie\_id INT REFERENCES Movie(movie\_id), artist\_id INT REFERENCES Artist(artist\_id), PRIMARY KEY(movie\_id, artist\_id) );

**Rating:** CREATE TABLE Rating ( rating\_id INT PRIMARY KEY,

);

movie\_id INT REFERENCES Movie(movie\_id) NOT NULL,
rating NUMBER(1, 1) NOT NULL
);

#### **PostProductionDoneIn:**

CREATE TABLE PostProductionDoneIn ( movie\_id INT REFERENCES Movie(movie\_id), building\_id INT REFERENCES Building(building\_id), PRIMARY KEY(movie\_id, building\_id) );

#### **BuildingPurpose:**

CREATE TABLE BuildingPurpose ( building\_id INT REFERENCES Building(building\_id), purpose VARCHAR2(100) NOT NULL, PRIMARY KEY(building\_id, purpose) );

#### **BuildingLocation:**

CREATE TABLE BuildingLocation ( building\_id INT REFERENCES Building(building\_id), location\_id INT REFERENCES SiteLocation(location\_id), PRIMARY KEY(building\_id, location\_id) );

The above tables satisfy the BCNF and 3NF requirements and all functional dependencies.

In summary, we identified 21 non-trivial functional dependencies among the original tables and normalized them into 22 tables that satisfy BCNF and 3NF requirements.

# The final Database:

#### Movie\_Script\_Inventory:

CREATE TABLE MovieScriptInventory ( script\_inventory\_id INT PRIMARY KEY, script\_inventory\_name VARCHAR(255) NOT NULL );



🔄 SQL Plus X + V	-	٥	×
SQL> select * from MovieScriptInventory;			
SCRIPT_INVENTORY_ID			
SCRIPT_INVENTORY_NAME			
The Shawshank Inventory			
2 The Avengers Inventory			
3 The Dark Knight Inventory			
SCRIPT_INVENTORY_ID			
SCRIPT_INVENTORY_NAME			
4 Fiction Inventory			
5 The Lord of the Rings Inventory			
6 Forrest Inventory			
SCRIPT_INVENTORY_ID			
SCRIPT_INVENTORY_NAME			
7			

#### Movies:

CREATE TABLE Movie (

movie\_id INT PRIMARY KEY,

title VARCHAR2(100) NOT NULL,

date\_of\_release DATE NOT NULL,

duration INT NOT NULL,

script\_inventory\_id INT REFERENCES MovieScriptInventory(script\_inventory\_id) NOT
NULL,

director VARCHAR2(100) NOT NULL,

in\_production CHAR(1) NOT NULL

);

SQL Plus



I SQL Plus	- 0	ζ
SQL> select * from MovieScriptInventory;		^
SCRIPT_INVENTORY_ID		
SCRIPT_INVENTORY_NAME		
1 The Shawshank Inventory		
The Avengers Inventory		
3 The Dark Knight Inventory		
SCRIPT_INVENTORY_ID		
SCRIPT_INVENTORY_NAME		
4 Fiction Inventory		
5 The Lord of the Rings Inventory		
Forrest Inventory		
SCRIPT_INVENTORY_ID SCRIPT_INVENTORY_NAME		
7 The Silence Inventory		
8 Matrix Inventory		
9 Ryan Inventory		~

#### **Movie Rating:**

CREATE TABLE MovieRating ( movie\_id INT PRIMARY KEY REFERENCES Movie(movie\_id), avg\_rating NUMBER(3, 1) NOT NULL );

```
SQL> CREATE TABLE MovieRating (
           movie_id INT PRIMARY KEY REFERENCES Movie(movie_id),
avg_rating NUMBER(3, 1) NOT NULL
  2
  З
  4
           );
CREATE TABLE MovieRating (
SQL> select * from MovieRating;
 MOVIE_ID AVG_RATING
       1 8.2
                7.5
        4
               6.9
                8.7
        6
                9.5
        8
                6.5
                 8.8
       10
                 7.3
10 rows selected.
SQL>
```

## Songs:

CREATE TABLE Songs ( song\_id INT NOT NULL, song\_name VARCHAR(255) NOT NULL, singer\_name VARCHAR(255) NOT NULL, movie\_id INT NOT NULL, PRIMARY KEY (song\_id), FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id) );

SQL> DROP TABLE Songs CASCADE CONSTRAINTS;	
Table dropped.	
SQL> rem +	
SQL> SQL> CREATE TABLE Songs ( song.in INT NOT NULL, song.name VARCHAR(255) NOT NULL, singer_name VARCHAR(255) NOT NULL, singer_name VARCHAR(255) NOT NULL, songer_name VARCHAR(255) NOT NULL,	
Table created. SQL>	
S SQL Plus X + V	- a x
SQL> select * from Songs;	
SONG_ID	
SONG_NAME	
SINGER_NAME	
MOVIE_ID	
1 Shape of You Ed Sheeran 1	
SONG_ID	
SONG_NAME	
SINGER_NAME	
MOVIE_ID	
2 Billie Jean Michael Jackson 2	
SINGER_NAME	

## Genre:

CREATE TABLE Genre ( genre\_id INT NOT NULL, genre\_name VARCHAR(255) NOT NULL,

# PRIMARY KEY (genre\_id)

);

SQL> DROP TABLE Genre CASCADE CONSTRAINTS;	
Table dropped.	
SQL> rem   Create Genre Table   SQL> rem   Create Genre Table   SQL> rem +	
Table created.	
sqL>	
🗟 SQL Plus X + V	×
SQL> select * from Genre;	
GENRE_ID	
GENRE_NAME	
1 Action	
2 Conedy	
3 Drama	
GENRE_ID	
GENRE_NAME	
4 Romance	
5 Thriller	
6 Adventure	
GENRE_ID	
GENRE_NAME	
7	

## **Tagged with:**

CREATE TABLE TaggedWith ( movie\_id INT NOT NULL, genre\_id INT NOT NULL, PRIMARY KEY (movie\_id, genre\_id), FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id), FOREIGN KEY (genre\_id) REFERENCES Genre(genre\_id)





SQL> select	* from Tag	jgedWith;
MOVIE_ID	GENRE_ID	
1	4	
1	3	
20		
2		
3		
3		
12		
5		
5		
13	9	
6	1	
MOVIE_ID	GENRE_ID	
	c	
11	1	
7	5	
8	6	
8	7	
13	8	
9		
9		
14		
10		
10		
22 rows sele	ected.	

## Sitelocation:

CREATE TABLE SiteLocation ( location\_id INT NOT NULL, location\_name VARCHAR(255) NOT NULL, address VARCHAR(255) NOT NULL, PRIMARY KEY (location\_id) );

SQL Plus		
SQL> select * from	Building;	
BUILDING_ID		
BUILDING_NAME		
PURPOSE	LOCATION_ID	
1		
MGM Studios studio	1	
2 Sony Pictures studio	2	
BUILDING_ID		
BUILDING_NAME		
PURPOSE	LOCATION_ID	
Pinewood Studios studio		
4 Warner Bros. Studio	5	
BUILDING_ID		

#### **Building:**

CREATE TABLE Building ( building\_id INT PRIMARY KEY, building\_name VARCHAR2(100) NOT NULL );

```
SQL> CREATE TABLE Building (

2 building_id INT PRIMARY KEY,

3 building_name VARCHAR2(100) NOT NULL

4 );

CREATE TABLE Building (
```

SQL> select * from Building;
BUILDING_ID
BUILDING_NAME
1 Empire State Building
2 Burj Khalifa
3 Taipei 101
BUILDING_ID
BUILDING_NAME
4 Shanghai Tower
5 The Shard
6 CN Tower
BUILDING_ID
BUILDING_NAME
7 Petronas Towers
8 Eiffel Tower

## ShotAt:

CREATE TABLE ShotAt (

location\_id INT NOT NULL,

movie\_id INT NOT NULL,

PRIMARY KEY (location\_id, movie\_id),

FOREIGN KEY (location\_id) REFERENCES SiteLocation(location\_id),

FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id)

);

SQL> DROP TABLE ShotAt CASCADE CONSTRAINTS;
Fable dropped.
SQL> rem
SQL> select * from ShotAt;
OCATAOLID MOVIE_ID
2 2
4 3
4 4 5 5
7 7 7 8
, , , , , , , , , , , , , , , , , , ,
LOCATION_ID HOVIE_ID
8 10
9   11   10   12
10 13
10 14 10 15
10 16
10 17 10 18
10 19
21 rows selected.
3QL>

#### **PostProductionDoneIn:**

CREATE TABLE PostProductionDoneIn ( movie\_id INT NOT NULL, building\_id INT NOT NULL, PRIMARY KEY (movie\_id, building\_id), FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id), FOREIGN KEY (building\_id) REFERENCES Building(building\_id) );

#### 

SQL> select	t * from Post	ProductionDoneIn;
MOVIE_ID	BUILDING_ID	
1	1	
2	2	
3	3	
4		
5		
6	6	
7		
8	8	
9	9	
10	10	
11	11	
MOVIE_ID	BUILDING_ID	
12	12	
13	13	
14	14	
15	15	
16	16	
17	17	
18	18	
19	19	
20	20	
20 rows se	lected.	
SQL>		

## **Employees:**

CREATE TABLE Employees (

employee\_id INT NOT NULL,

employee\_name VARCHAR(255) NOT NULL,

designation VARCHAR(255) NOT NULL,

phone\_number VARCHAR(10) NOT NULL UNIQUE,

PRIMARY KEY (employee\_id),

CONSTRAINT designatoin\_constraint\_violated CHECK (

designation IN ('choreographers', 'security', 'sound engineer', 'makeup artist', 'electrician', 'janitor', 'manager')));



SQL Plus		
SQL> select * from I	Employees;	
EMPLOYEE_ID		
EMPLOYEE_NAME		
DESIGNATION		
PHONE_NUMB		
101		
John Doe		
electrician		
1234567890		
EMPLOYEE_ID		
EMPLOYEE_NAME		
DESIGNATION		
PHONE_NUMB		
102		
Jane Smith		
7345678901		
EMPLOYEE TD		
EMPLOYEE_NAME		
DESIGNATION		

## Payroll:

CREATE TABLE Payroll (

payroll\_id INT NOT NULL,

salary FLOAT NOT NULL,

employee\_id INT NOT NULL,

hours\_worked FLOAT NOT NULL,

PRIMARY KEY (payroll\_id),

FOREIGN KEY (employee\_id) REFERENCES Employees(employee\_id),

CONSTRAINT salary\_constraint\_violated CHECK (salary > 0),

CONSTRAINT hours\_constraint\_violated CHECK (hours\_worked > 0 and hours\_worked <= 40)

);

SQL> DROP TABL	E Payroll (	ASCADE CONSTRA	AINTS;
Table dropped.			
SOL> rem +			
SQL> rem   Cre	ate Payroll	. Table	
SQL> rem + SQL>			
SQL> CREATE TA	BLE Payroll	. (	
2 payro 3 salar	V FLOAT NOT	IOT NULL, I NULL.	
4 emplo	yee_id INT	NOT NULL,	
5 hours	_worked FLC	DAT NOT NULL,	
7 FOREI	GN KEY (emp	oloyee_id) REF	ERENCES Employee
8 CONST	RAINT salar	y_constraint_	violated CHECK (
9 CONST	RAINT hours	<pre>s_constraint_v:</pre>	iolated CHECK (h
10 ),			
Table created.			
SOL>			
SQL> select	* from Pa	yroll;	
	CAL ADV		
	JALART	ENPLOYEE_ID	
1	5000	101	40
	6000	102	38.5
	4500	103	37
4	5500	104	39.5
5	4000	105	36
5	4800	105	40 38.5
8	5200	107	37
9	5800	109	39.5
10	4200	110	36
11	5100	111	40
PAYROLL_ID	SALARY	EMPLOYEE_ID	HOURS_WORKED
12	5300	112	38.5
13	4700	112	37
14	5400	114	39.5
15	3900	115	36
16	5900	116	40
17	4600	117	38.5
18	5000	118	37
29	4300	119	39.5
20	4500	120	
20 rows sele	cted.		
SQL>			

## **Sponsoring company**

CREATE TABLE SponsoringCompany ( company\_id INT NOT NULL, company\_name VARCHAR(255) NOT NULL UNIQUE, PRIMARY KEY (company\_id)

);



🗟 SQL Plus X + V	-	ð	×
SQL> select * from SponsoringCompany ;			
COMPANY_ID			
COMPANY_NAME			
1 Universal Pictures			
2 Warner Bros. Pictures			
3 Walt Disney Pictures			
COMPANY_ID			
COMPANY_NAME			
4 Paramount Pictures			
5 20th Century Fox			
6 Sony Pictures			
COMPANY_ID			
COMPANY_NAME			
7			

#### Manages:

CREATE TABLE Manages ( employee\_id INT NOT NULL, location\_id INT NOT NULL, PRIMARY KEY (employee\_id, location\_id), FOREIGN KEY (employee\_id) REFERENCES Employees(employee\_id), FOREIGN KEY (location\_id) REFERENCES SiteLocation(location\_id) );



SQL> select	* from Manage	es;	
EMPLOYEE_ID	LOCATION_ID		
101	1		
102	2		
103			
104	4		
105			
106	6		
107	7		
108	8		
109	9		
110	10		
101	2		
EMPLOYEE_ID	LOCATION_ID		
102	3		
103	4		
104	5		
105	6		
106			
107	8		
108	9		
109	10		
110	1		
20 rows sele	cted.		
sql>			

#### Artist:

CREATE TABLE Artist ( artist\_id INT PRIMARY KEY, artist\_name VARCHAR(255) NOT NULL, date\_of\_birth DATE NOT NULL, gender VARCHAR(1) NOT NULL, CONSTRAINT gender\_constraint\_violated CHECK (gender in ('M', 'F', 'T')) );



SQL Plus	x + •	- a	×
SQL> select * from Artis	st;		
ARTIST_ID			
ARTIST_NAME			
DATE_OF_B G			
Tom Hanks 09-JUL-56 M			
2 Meryl Streep 22-JUN-49 F			
ARTIST_ID			
ARTIST_NAME			
DATE_OF_B G			
3 Denzel Washington 28-DEC-54 M			ſ
4 Scarlett Johansson			
ARTIST_ID			
ARTIST_NAME			
DATE_OF_B G			

## GetsPaidBy

CREATE TABLE getsPaidBy ( artist\_id INT NOT NULL, company\_id INT NOT NULL, PRIMARY KEY (artist\_id, company\_id), FOREIGN KEY (artist\_id) REFERENCES Artist(artist\_id), FOREIGN KEY (company\_id) REFERENCES SponsoringCompany(company\_id) );

SOLS DROD TARLE RateDaidRy CASCADE CONSTDATIONS						
see blor indet gestatug chsche considering,						
Table dropped.						
SQL> rem +       +         SQL> rem +       -         SQL> rem +       -         SQL> rem +       -         SQL> CREATE TABLE getsPaidBy (       -         SQL> CREATE TABLE getsPaidBy (       2         artist_id INT NOT MULL,       -						
3 company_id INT NOT NULL, 4 PRIMARY KEY (artist_id, company_id), 5 FORFOR WEY (artist_id) DEFENSIVES						
5 FURLIN REY (artist_10) REFERENCES Artist(artist_10), 6 FOREIGN REY (company_id) REFERENCES SponsoringCompany_id) 7 );						
Table created.						
SQL>						
SQL> select * from getsPaidBy;						
ARTIST_ID COMPANY_ID						
1 4						
2 4						
2 5						
3 3						
3 4 4 1						
4 2						
ARTIST_ID COMPANY_ID						
4 3						
5 1 5 2						
15 rows selected.						
SOL>						

Produces: CREATE TABLE Produces ( company\_id INT NOT NULL, movie\_id INT NOT NULL, PRIMARY KEY (company\_id, movie\_id), FOREIGN KEY (company\_id) REFERENCES SponsoringCompany(company\_id), FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id) );



#### ActsIn:

CREATE TABLE ActsIn ( movie\_id INT NOT NULL, artist\_id INT NOT NULL, PRIMARY KEY (movie\_id, artist\_id), FOREIGN KEY (movie\_id) REFERENCES Movie(movie\_id), FOREIGN KEY (artist\_id) REFERENCES Artist(artist\_id) );

SQL> DROP TABLE ACTSIN CASCADE CONSTRAINTS;							
Table dropped.							
SQL> rem +	I Table   IT NULL, IT NULL, IOT NULL, Vie_id, artist_id), Vie_id) REFERENCES Movie(movie_id), tist_id) REFERENCES Artist(artist_id)						
Table created.							
SQL>		1					
SQL> select * from Ac	stsIn;						
MOVIE_ID ARTIST_ID							
1 1 2 2 3 3 4 4 5 5 6 6 7 7 8 8 9 9 10 16 5 12 15 0 12							
7 14 19 13 2 12 15 13 4 16 8 11 18 12 13 12 5 12							
20 rows selected.							
sql>							

# **Individual Contribution:**

Identified Functional Dependencies in Movie and Songs table.

Designed Movie and Songs tables in 2NF.

Modified Movie and Songs tables in 3NF.

Modified Movie and Songs tables to BCNF.

Found Dependency preserving and loss less join in Movie and Songs tables.

Created table for Movie and Songs.